

# Counteracting UDP Flooding Attacks in SDN

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# Outline

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- **SDN** overview
- Problem statement
- Proposed method
- Experiments

# SDN Introduction

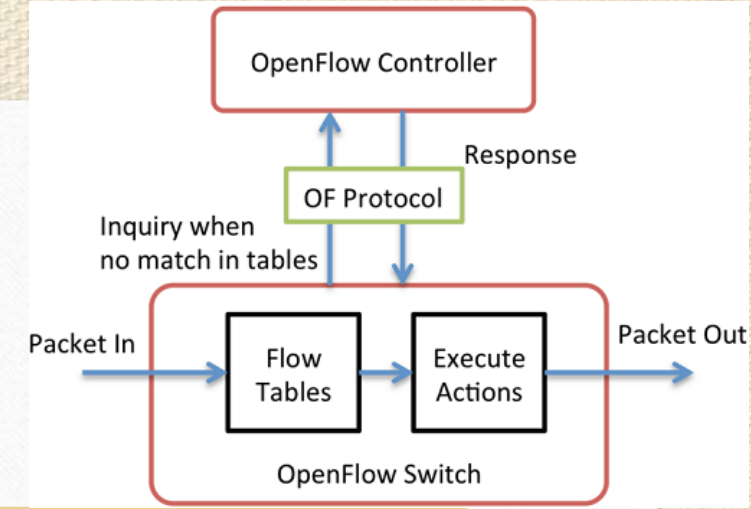
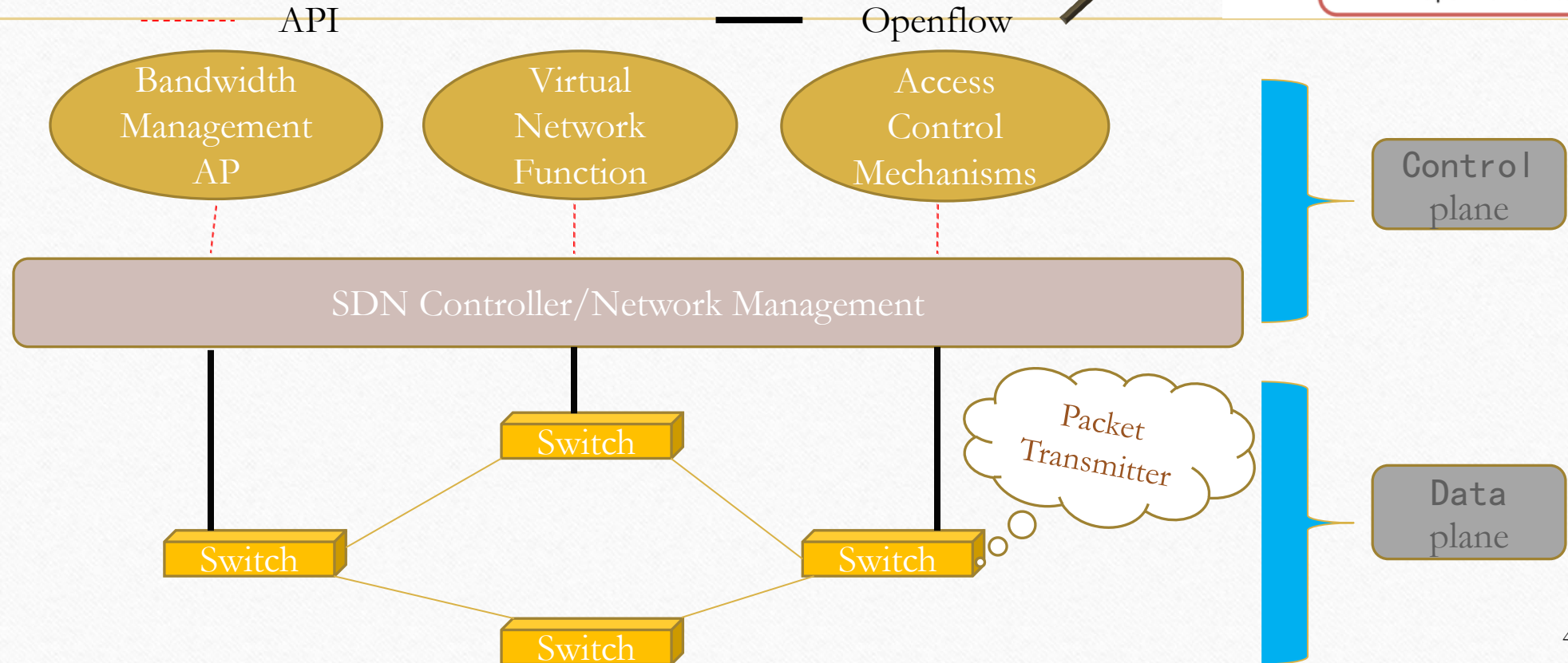
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- Centralized approach
- SDN mainly divided into control plane and data plane
- SDN uses the OpenFlow protocol
- SDN switch has a flow table, trying to have a rule match against the received packets



# SDN Introduction

## Framework:



# Problem Statement

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- **Network Security**

- The easiest way of compromising a network is to launch a flooding attack (ex: TCP SYN flooding, UDP flooding etc ).

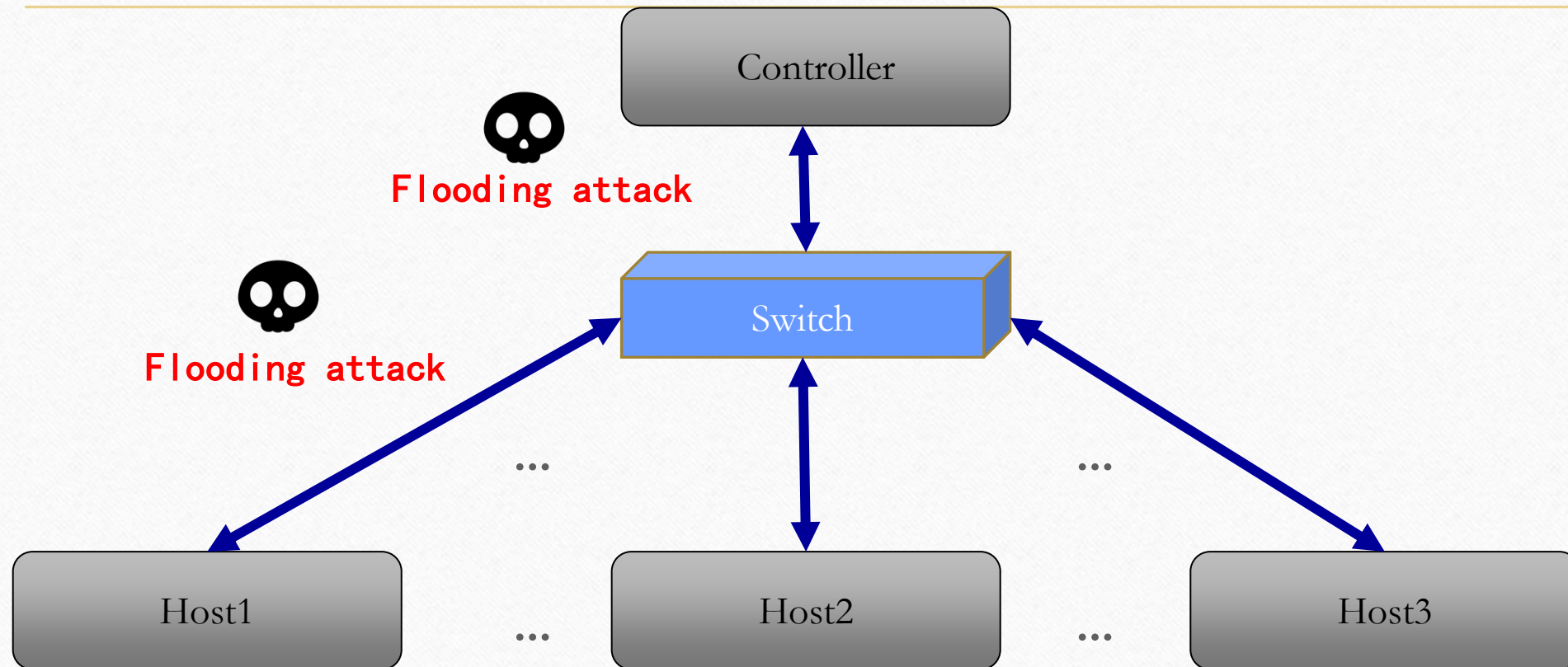
- **SDN Security Problems**

- When a new flow arrives, the SDN switch will send a packet-in message to the SDN controller.
- However, intentional abusing the controller (or say packet-in message) may incur the security problem.



# Problem Statement

## Simulation SDN Network Attack Graph



# PROTOCOL DESIGN

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- Our experiment can be divided into two phases
  - First, consider a bunch of simple UDP packets transmitted to the switch.
  - Then, we began to do the code implementation on the simulated switch and controller, and evaluated the performance and the security of our defense mechanism.



# PROTOCOL DESIGN

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## Attack Model:

- In the case of no match found, the controller will perform a broadcast to ask whether there is a match for the purpose of IP addresses.
- The attacker can assign a random value to the destination field in the packet.

```
def generate_ip(): # Create random IP
    return str(random.randint(0, 255)) + '.' + str(random.randint(0, 255)) + '.' + str(random.randint(0, 255)) + '.' + str(random.randint(0, 255))
```



# PROTOCOL DESIGN

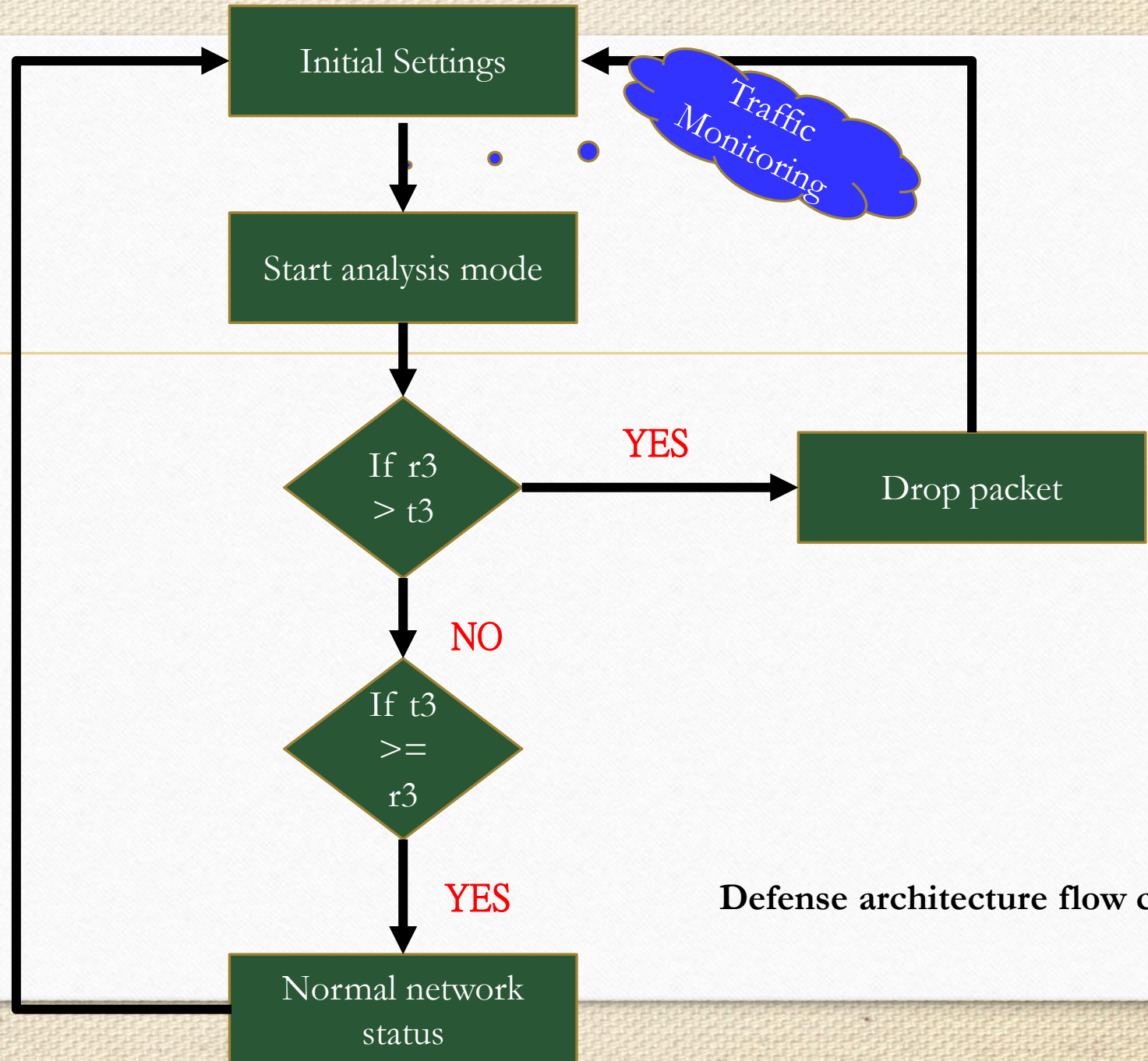
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	Total Rate	CPU (s)	Load avg.
<b>Normal state</b>	5 kbits/sec	0.6 us	0.32
<b>Attack state</b>	6100 ↑ kbits/sec	27 ↑ us	0.87 ↑

# Defense Architecture

r3 : The number of packets receive by the port

t3 : The number of packets sent by the port



Defense architecture flow chart



# Defense Architecture

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- Our analysis model has two conditions.
  - If the received packet (r3)  $>$  send packets (t3):
    - This means that the destination of the sending packet does not exist in the current network, resulting in the controller constantly broadcasting.
  - If the packet is sent (t3)  $> =$  receive packets (r3):
    - The controller can handle the packet-in message and broadcast packets.

# Defense Architecture

## UDP Defense Section

```
@set_ev_cls(ofp_event.EventOFPPacketIn, MAIN_DISPATCHER)
def _packet_in_handler(self, ev):
    if ev.msg.msg_len < ev.msg.total_len: self.logger.debug("packet truncated: only %s of %s bytes",
        ev.msg.msg_len, ev.msg.total_len)
    .
    .
    if(r3 > t3):
        actions = []
    .
    .
    elif(t3 >= r3):
        flooding
```



# Defense Architecture

## Return packets on all ports

```
body = ev.msg.body
self.logger.info('datapath      port      '
                 'rx-pkts rx-bytes rx-error '
                 'tx-pkts tx-bytes tx-error')
self.logger.info('----- '
                 '----- '
                 '-----')
for stat in sorted(body, key=attrgetter('port_no')):
self.logger.info('%016x %08x %08d %08d %08d %08d %08d %08d',
                 ev.msg.datapath.id, stat.port_no,
                 stat.rx_packets, stat.rx_bytes, stat.rx_errors,
                 stat.tx_packets, stat.tx_bytes, stat.tx_errors)
```

# EXPERIMENTS

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- Experiment Setting
  - In the experiment. we use `mininet` to simulate the SDN `OpenFlow` switch, and use `RYU` to simulate the controller.
  - Moreover, `IPerf`, `TOP`, `IPTRAF` are used as monitoring tools.
  - For the network topology, we considered two physical hosts and a controller.
  - They are on different physical machines for ensuring more accurate measurement.



# EXPERIMENTS

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- **Defense Achievements**

- In our experiment, we consider two cases (with and without attack) and observe the difference between these two cases.

# EXPERIMENTS

## Network bandwidth and controller performance comparison

	IPerf	Top	IPtraf
No Defense	TX bps:412 Bytes/s	CPU(s) : 27.2 us	Total rate: 6139.0 Kbits/sec 4846.4 packets/sec
Defense	TX bps: 33 Bytes/s	CPU(s) : 14.8 us	Total rate: 2790.7 Kbits/sec 1861.8 packets/sec



# Related Work

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- Comparison of Defense

	FloodGuard	UDP
No Defense	7 Mbps	6 Mbps
Defense	2 Mbps	2 Mbps

# CONCLUSION

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- The proposed defense resist against the UDP flooding with a minor modification in SDN module.
- The countermeasure particularly designed for only UDP flooding works with better performance



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Let us know if you have any comments or questions.

Thank you for listening.

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# Question

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- Given the operation flow chart probing the switches periodically, it would be a naturally raised question how much overhead this approach would introduce.
- Furthermore, this question extends to what is the parameters we should consider to trade off security and performance compromise.



# Answer

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- Using this method, we are only at the expense of request packet for some time. The following mechanisms to facilitate the analysis.
- Although this sacrifices some benign request, but in exchange for increased security.
- But in the time of the attack, a benign request to wait for a short time.

# Question

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- The conditions, ' $r_3 > t_3$ ' or ' $t_3 \geq r_3$ ' over simplifies or ignores lots of other possibilities considering the nature of UDP traffic ( eg. streaming applications).



# Answer

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- Perhaps while watching the movie, the flow slightly. But the normal traffic.
- This time we use to calculate packet per second to reduce false positives.